

# IM 7011: Information Economics

Lecture 12: Moral Hazard  
Chen and Huang (2013)

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November 25, 2013

# Road map

- ▶ **Introduction.**
- ▶ Simplified model.
- ▶ Analysis.
- ▶ Original model and analysis.
- ▶ Extensions and conclusions.

## Pricing data services

- ▶ We use **data services** everyday.
  - ▶ Text messages.
  - ▶ Dial-up or ADSL.
  - ▶ 3G/4G.
- ▶ How do sellers (e.g., ISPs) **price** these services?
  - ▶ Text messages: by **quantity**.
  - ▶ Dial-up: by **time**.
  - ▶ ADSL: by **bandwidth**.
  - ▶ 3G/4G: by volume (i.e., quantity).
- ▶ Why different data services are priced by different **pricing metrics**?
  - ▶ There are certainly **supply-side** reasons, e.g., technology limits.
  - ▶ Is there any **consumer-side** reasons?
- ▶ Practitioners often make (effective or ineffective) decisions without using **scientific** methods.
  - ▶ We want to know whether pricing metrics are chosen in a “good” way.

## Pricing metrics

- ▶ Suppose a monopoly data service provider (**seller**) intends to provide the services to **consumers**.
  - ▶ In the basic model, the cost for offering services are omitted.
  - ▶ The seller wants to find the **revenue-maximizing pricing** plan.
- ▶ Consumers are **heterogeneous** on their willingness-to-pay for data usage and connection speed.
- ▶ As consumer types are hidden, the seller can only adopt second- or third-degree price discrimination.<sup>1</sup>
- ▶ We will focus on second-degree price discrimination with the following three **pricing metrics**:
  - ▶ Pricing by time (e.g., minutes).
  - ▶ Pricing by bandwidth (e.g., Mbps).
  - ▶ Pricing by quantity (e.g., Gigs).
- ▶ Which pricing metric is the best?

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<sup>1</sup>Pricing by usage/choice or attribute/identity.

## After-sales selections

- ▶ Consumers do not just have hidden types.
- ▶ They also have hidden (uncontrolled) **after-sales selections**.
  - ▶ When I am priced by time, I select connection speed (by selecting software/applications).
  - ▶ When I am priced by bandwidth, I select my time usage.
  - ▶ When I am priced by quantity, I select time or speed.
- ▶ Each consumer acts to maximize his own utility.
- ▶ The selection of pricing metrics must consider:
  - ▶ The heterogeneity of consumers (**hidden information**).
  - ▶ The after-sales selections (**hidden action**).

## Research questions

- ▶ The seller wants to find the **revenue-maximizing pricing metric**.
  - ▶ By time, bandwidth, or quantity?
- ▶ To answer this question, she must be able to find the optimal (second-best) menu under each pricing metric.
  - ▶ Given each pricing metric, the seller solves a **nonlinear pricing problem** through contract design.
  - ▶ Multi-tiered pricing, unlimited usage pricing, or both?
- ▶ To solve the nonlinear pricing problem, the seller must be able to **anticipate** each consumers' after-sales selection.
- ▶ As researchers, we want to find the **driving forces** for a pricing metric to be revenue-maximizing.
  - ▶ When one is better than the other, and **why**?

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## Pricing metrics

- ▶ A monopoly risk-neutral seller is facing three options:
  - ▶ Pricing by minutes ( $M$ ).
  - ▶ Pricing by bandwidth ( $B$ ).
  - ▶ Pricing by quantity ( $Q \equiv BM$ ).
- ▶ For pricing by  $M$  and  $Q$ , we exclude fixed-up-to plans.
  - ▶ Fixed-up-to plans may arise as a consequence of optimization.
  - ▶ We do not specifically focus on such a restriction.
- ▶ Given a pricing metric, the seller designs a **price schedule**.
  - ▶ For example, under pricing by minutes, the seller designs a function  $P^M(M)$  to translate a time usage  $M$  to a payment  $P^M(M)$ .
- ▶ A price schedule can be implemented as a **menu of contracts**.
  - ▶ For example,  $P^M(\cdot)$  can be implemented as  $\{(M(\theta), P^M(\theta))\}$ , where  $\theta$  is the consumer's type (to be detailed later).
  - ▶ A price schedule is an indirect mechanism; a menu is a direct one.



## Consumers' utility function

- ▶ Let  $\theta \sim \text{Uni}(0, 1)$  be the consumers' type.
- ▶ In the **simplified** model,<sup>2</sup> the type- $\theta$  consumer's utility is<sup>3</sup>

$$u(B, M, \theta) = \begin{cases} \theta BM - \frac{1}{2}(BM)^2 & + \theta B - \frac{1}{2}B^2 & \text{if } BM \leq \theta \text{ and } B \leq \theta \\ \frac{1}{2}\theta^2 & + \theta B - \frac{1}{2}B^2 & \text{if } BM > \theta \text{ and } B \leq \theta \\ \theta BM - \frac{1}{2}(BM)^2 & + \frac{1}{2}\theta^2 & \text{if } BM \leq \theta \text{ and } B > \theta \\ \frac{1}{2}\theta^2 & + \frac{1}{2}\theta^2 & \text{if } BM > \theta \text{ and } B > \theta \end{cases} .$$

- ▶ The first part ( $\theta BM - \frac{1}{2}(BM)^2$  and  $\frac{1}{2}\theta^2$ ) makes  $u(\cdot)$  **increasing and concave** in  $Q$ .
- ▶ They also make  $u(\cdot)$  increasing and concave in  $M$  when  $B$  is fixed.
- ▶ The second part ( $\theta B - \frac{1}{2}B^2$  and  $\frac{1}{2}\theta^2$ ) makes  $u(\cdot)$  increasing and concave in  $B$  when  $Q$  is fixed.
- ▶ **Unlimited usage** does not give unlimited utility.

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<sup>2</sup>We remove some parameters from the paper's original model at this moment.

<sup>3</sup>The "if" condition in the paper should be a typo. The sign should be reversed.

## More about consumers' utility function

- ▶ The functional form

$$\theta BM - \frac{1}{2}(BM)^2 + \theta B - \frac{1}{2}B^2$$

has its limitations.

- ▶ Consumers who have stronger preference for  $Q$  **also** have stronger preference for  $B$ .
- ▶ Nevertheless, multi-dimensional screening is too hard.
- ▶ A higher time usage results in a higher utility **only if** it corresponds to a higher data usage.
  - ▶ Consuming more time itself does not make one happier.
- ▶ As there is no cost for offering the service, the **socially efficient** consumption maximizes each consumer's utility.
  - ▶ The FOC gives  $B = \frac{\theta(1+M)}{1+M^2}$  and  $M = \frac{\theta}{B}$ , which imply  $B = \theta$  and  $M = 1$ .
  - ▶ Will there be efficiency loss?

# Timing

- ▶ The seller determines the pricing metric.
- ▶ The seller announces a pricing menu.
  - ▶ For example, if she prices by minutes, she announces  $\{(M(\theta), P^M(\theta))\}$ .
- ▶ Each consumer self-selects one contract in the menu.
- ▶ Each consumer adjusts the variable not specified in the contract.
  - ▶ For example, if the seller prices by minutes, the consumer chooses his own connection speed.

# Road map

- ▶ Introduction.
- ▶ Simplified model.
- ▶ **Analysis.**
  - ▶ Pricing by minutes.
  - ▶ Pricing by bandwidth.
  - ▶ Pricing by quantity.
  - ▶ Comparisons.
- ▶ Original model and analysis.
- ▶ Extensions and conclusions.

## Pricing by minutes: after-sales selection

- ▶ Suppose the type- $\theta$  consumer has chosen  $(M(\hat{\theta}), P^M(\hat{\theta}))$  in stage 3.
- ▶ In stage 4, he determines the bandwidth  $B$  to maximize his **net utility**

$$U^M(B|\theta, \hat{\theta}) = \theta BM(\hat{\theta}) - \frac{1}{2} \left( BM(\hat{\theta}) \right)^2 + \theta B - \frac{1}{2} B^2 - P^M(\hat{\theta}).$$

- ▶ To maximize his net utility, the consumer chooses the bandwidth

$$B^*(\theta, \hat{\theta}) = \theta \left[ \frac{1 + M(\hat{\theta})}{1 + M(\hat{\theta})^2} \right].$$

- ▶ The **effective utility** of choosing  $(M(\hat{\theta}), P^M(\hat{\theta}))$  is

$$U^M(\theta, \hat{\theta}) = \frac{\theta^2}{2} \frac{[1 + M(\hat{\theta})]^2}{1 + M(\hat{\theta})^2} - P^M(\hat{\theta}).$$

- ▶ Let  $U^M(\theta) \equiv \max \{ U^M(\theta, \theta), 0 \} \equiv [U^M(\theta, \theta)]^+$ .

## Pricing by minutes: contract design

- ▶ In stage 2, the seller solves

$$\begin{aligned}\Pi^M &= \max_{M(\cdot), P^M(\cdot)} \mathbb{E} \left[ P^M(\theta) \right] \\ \text{s.t. } & U^M(\theta) \geq U^M(\theta, \hat{\theta}) \quad \forall \theta, \hat{\theta} \\ & U^M(\theta) \geq 0 \quad \forall \theta.\end{aligned}$$

- ▶ To solve this problem, we apply the standard technique for continuous-type problems and other recent results.

## Pricing by minutes: optimal menu

- ▶ It turns out that a **fixed-fee** pricing plan is optimal.

### Lemma 1

*Under pricing by minutes, the optimal pricing plan is to charge a single fixed fee  $P^M = \frac{4}{9}$  for an unlimited usage. The seller's expected revenue is  $\Pi^M = \frac{4}{27}$ .*

- ▶ By buying the unlimited time usage, the type- $\theta$  consumer's net utility becomes

$$\frac{1}{2}\theta^2 + \frac{1}{2}\theta^2 - P^M.$$

Therefore, he buys the service if and only if  $\theta \geq \sqrt{P^M}$ .

- ▶ The seller then maximizes the expected revenue  $P^M(1 - \sqrt{P^M})$ .
- ▶ Price discrimination is **suboptimal**.
- ▶ In equilibrium the seller **does not screen** consumers!

## Pricing by bandwidth: after-sales selection

- ▶ Suppose the type- $\theta$  consumer has chosen  $(B(\hat{\theta}), P^B(\hat{\theta}))$  in stage 3.
- ▶ In stage 4, he determines the time usage  $M$  to maximize

$$U^B(M|\theta, \hat{\theta}) = \theta B(\hat{\theta})M - \frac{1}{2} [B(\hat{\theta})M]^2 + B(\hat{\theta})\theta - \frac{1}{2} B(\hat{\theta})^2 - P^B(\hat{\theta}).$$

- ▶  $M$  only appears in the first part (quantity).
- ▶ The consumer chooses the time usage  $M^*(\theta, \hat{\theta}) = \frac{\theta}{B(\hat{\theta})}$ .
- ▶ The effective utility of choosing  $(B(\hat{\theta}), P^B(\hat{\theta}))$  is

$$U^B(\theta, \hat{\theta}) = \frac{1}{2}\theta^2 + B(\hat{\theta})\theta - \frac{1}{2}B(\hat{\theta})^2 - P^B(\hat{\theta}).$$

- ▶ Let  $U^B(\theta) \equiv \max \{U^B(\theta, \theta), 0\} \equiv [U^B(\theta, \theta)]^+$ .



## Pricing by bandwidth: contract design

- ▶ In stage 2, the seller solves

$$\begin{aligned}\Pi^B &= \max_{B(\cdot), P^B(\cdot)} \mathbb{E}\left[P^B(\theta)\right] \\ \text{s.t. } & U^B(\theta) \geq U^B(\theta, \hat{\theta}) \quad \forall \theta, \hat{\theta} \\ & U^B(\theta) \geq 0 \quad \forall \theta.\end{aligned}$$

## Pricing by bandwidth: optimal menu

- ▶ Now **multi-tiered** (usage-based) pricing is optimal.

### Lemma 2

*Under pricing by bandwidth, the optimal pricing plan satisfies*

$$B^*(\theta) = 2\theta - 1 \quad \text{and} \quad P^B(\theta) = 2\theta - \theta^2 - \frac{1}{2} + \frac{\theta(2\underline{\theta}^2 - \underline{\theta} + 3)}{2(3\underline{\theta} - 2)}$$

*for  $\theta \geq \underline{\theta}$  and  $B^*(\theta) = P^B(\theta) = 0$  for  $\theta < \underline{\theta}$ , where  $\underline{\theta} = \frac{3+\sqrt{2}}{7}$  is the lowest type of consumer that is served. The seller's expected revenue is  $\Pi^B = \frac{1}{6} - \underline{\theta}^2(\frac{3}{2} - \frac{7}{3}\underline{\theta})$ .*

- ▶ Monotonicity:  $B^*(\theta)$  is nondecreasing. Also no rent at bottom.
- ▶ Efficiency at top:  $B^*(\theta) = 2\theta - 1 = \theta \Leftrightarrow \theta = 1$ .
- ▶ Price discrimination is optimal but some consumers should be ignored.
- ▶ **Quantity discount**:  $B^*(\theta)$  is linear while  $P^B(\theta)$  is strictly concave.

## Pricing by quantity: after-sales selection

- ▶ Suppose the type- $\theta$  consumer has chosen  $(Q(\hat{\theta}), P^Q(\hat{\theta}))$  in stage 3.
- ▶ In stage 4, he determines the bandwidth  $B$  to maximize<sup>4</sup>

$$U^Q(B|\theta, \hat{\theta}) = \theta Q(\hat{\theta}) - \frac{1}{2}Q(\hat{\theta})^2 + B\theta - \frac{1}{2}B^2 - P^Q(\hat{\theta}).$$

- ▶  $B$  only appears in the second part (bandwidth).
- ▶ The consumer chooses the bandwidth  $B^*(\theta, \hat{\theta}) = \theta$ .
- ▶ The effective utility of choosing  $(B(\hat{\theta}), P^B(\hat{\theta}))$  is

$$U^Q(\theta, \hat{\theta}) = Q(\hat{\theta})\theta - \frac{1}{2}Q(\hat{\theta})^2 + \frac{1}{2}\theta^2 - P^Q(\hat{\theta}).$$

- ▶ Let  $U^Q(\theta) \equiv \max \{U^Q(\theta, \theta), 0\} \equiv [U^Q(\theta, \theta)]^+$ .

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<sup>4</sup>As long as  $Q(\hat{\theta}) = BM$ , an equivalent result may be obtained by using the time usage  $M$  as the variable or by using both  $B$  and  $M$  as variables.

## Pricing by quantity: contract design

- ▶ In stage 2, the seller solves

$$\begin{aligned}\Pi^Q &= \max_{Q(\cdot), P^Q(\cdot)} \mathbb{E}\left[P^Q(\theta)\right] \\ \text{s.t. } & U^Q(\theta) \geq U^Q(\theta, \hat{\theta}) \quad \forall \theta, \hat{\theta} \\ & U^Q(\theta) \geq 0 \quad \forall \theta.\end{aligned}$$

## Pricing by quantity: optimal menu

- ▶ Again, multi-tiered (usage-based) pricing is optimal.

### Lemma 3

*Under pricing by quantity, the optimal pricing plan satisfies*

$$Q^*(\theta) = 2\theta - 1 \quad \text{and} \quad P^Q(\theta) = 2\theta - \theta^2 - \frac{1}{2} + \frac{\theta(2\theta^2 - \theta + 3)}{2(3\theta - 2)}$$

*for  $\theta \geq \underline{\theta}$  and  $Q^*(\theta) = P^Q(\theta) = 0$  for  $\theta < \underline{\theta}$ , where  $\underline{\theta} = \frac{3+\sqrt{2}}{7}$  is the lowest type of consumer that is served. The seller's expected revenue is  $\Pi^Q = \frac{1}{6} - \underline{\theta}^2(\frac{3}{2} - \frac{7}{3}\underline{\theta})$ .*

- ▶ Identical to pricing by bandwidth!
- ▶ Consumers' effective utility is:
  - ▶  $\frac{1}{2}\theta^2 + B(\hat{\theta})\theta - \frac{1}{2}B(\hat{\theta})^2 - P^B(\hat{\theta})$  when pricing by bandwidth.
  - ▶  $Q(\hat{\theta})\theta - \frac{1}{2}Q(\hat{\theta})^2 + \frac{1}{2}\theta^2 - P^Q(\hat{\theta})$  when pricing by quantity.

## Selection among pricing metrics

- ▶ Now we may find the revenue-maximizing pricing metric:

### Proposition 1

- ▶ *A single contract is offered under pricing by minutes. A menu is offered under pricing by bandwidth or quantity.*
  - ▶ *Because  $\Pi^M \approx 0.148 < 0.155 \approx \Pi^B = \Pi^Q$ , pricing by minutes is not revenue-maximizing.*
  - ▶ *Because  $1 - \frac{2}{3} \approx 0.33 < 0.37 \approx 1 - \underline{\theta}$ , more consumers are served under pricing by bandwidth or quantity.*
  - ▶ *Pricing by bandwidth and pricing by quantity are equivalent.*
- ▶ Pricing by minutes **cannot screen** consumers (with a fixed fee).
  - ▶ Pricing by minutes is the least effective in **alleviating the moral hazard problem**.
    - ▶ Consumers are “too free”: They can adjust bandwidth to affect both bandwidth and quantity.
    - ▶ In the other two cases, only one part can be adjusted.

## Robustness of insights

- ▶ Are the insights robust?
  - ▶ Is pricing by minutes always inferior?
  - ▶ Are pricing by bandwidth and pricing by quantity always identical?
- ▶ To answer this question, a more general model is required.

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## Original consumers' utility function

- ▶ In the original model in the paper, the type- $\theta$  consumer's utility function is

$$u(B, M, \theta) = \begin{cases} \delta\theta BM - \frac{1}{2\eta}(BM)^2 & + \theta B - \frac{1}{2\gamma}B^2 & \text{if } BM \leq \theta \text{ and } B \leq \theta \\ \frac{1}{2}\eta\delta^2\theta^2 & + \theta B - \frac{1}{2\gamma}B^2 & \text{if } BM > \theta \text{ and } B \leq \theta \\ \delta\theta BM - \frac{1}{2\eta}(BM)^2 & + \frac{1}{2}\gamma\theta^2 & \text{if } BM \leq \theta \text{ and } B > \theta \\ \frac{1}{2}\eta\delta^2\theta^2 & + \frac{1}{2}\gamma\theta^2 & \text{if } BM > \theta \text{ and } B > \theta \end{cases} .$$

- ▶  $\delta > 1$  ( $\delta < 1$ ): One is more (less) **sensitive** to changes in  $Q$  than  $B$ .
  - ▶  $\eta$  ( $\gamma$ ) increases: The marginal benefit of quantity (bandwidth) diminishes in a **slower** rate.
- 
- ▶ With the more general utility function, do the results change?

## More general insights

- ▶ The old results can now be generalized:

### Proposition 2

- ▶ *A single contract is offered under pricing by minutes. A menu is offered under pricing by bandwidth or quantity.*
  - ▶ *Because  $\Pi^M < \Pi^B$  and  $\Pi^M < \Pi^Q$ , pricing by minutes is not revenue-maximizing.*
  - ▶ *Pricing by bandwidth is revenue-maximizing if and only if  $\gamma \geq \delta^2 \eta$ .*
- ▶ Some insights are robust:
    - ▶ Pricing by minutes still cannot screen consumers.
    - ▶ Pricing by minutes is still suboptimal.
  - ▶ Some are not:
    - ▶ Pricing by bandwidth and pricing by quantity are **not** identical.
    - ▶ **Both** of them may be revenue-maximizing.

## Revenue maximization and moral hazard

- ▶ Why pricing by bandwidth is optimal if and only if  $\gamma \geq \delta^2 \eta$ ?
- ▶ It depends on which pricing metric is more effective in alleviating the **moral hazard** issue.
  - ▶ Under pricing by bandwidth, the utility is

$$\underbrace{\delta \theta B(\hat{\theta})M - \frac{1}{2\eta} [B(\hat{\theta})M]^2}_{\text{can be adjusted}} + B(\hat{\theta})\theta - \frac{1}{2\gamma} B(\hat{\theta})^2.$$

- ▶ Under pricing by quantity, the utility is

$$\delta \theta Q(\hat{\theta}) - \frac{1}{2\eta} Q(\hat{\theta})^2 + \underbrace{B\theta - \frac{1}{2\gamma} B^2}_{\text{can be adjusted}}.$$

- ▶ When  $\gamma$  is large,  $B\theta - \frac{1}{2\gamma} B^2$  is large and pricing by quantity leaves the consumer **a too large room** for adjustment.
- ▶ When  $\delta$  or  $\eta$  is large,  $\delta \theta B(\hat{\theta})M - \frac{1}{2\eta} [B(\hat{\theta})M]^2$  is large.

## Revenue maximization and adverse selection

- ▶ Why pricing by bandwidth is optimal if and only if  $\gamma \geq \delta^2\eta$ ?
- ▶ It also depends on which pricing metric is more effective in alleviating the **adverse selection** issue.
- ▶ For the functional form

$$\delta\theta BM - \frac{1}{2\eta}(BM)^2 + \theta B - \frac{1}{2\gamma}B^2 :$$

- ▶ When  $\delta < 1$ , consumers are **more heterogeneous** in  $B$  than in  $Q$ .<sup>5</sup>
- ▶ Pricing by bandwidth, which screens consumers according to their willingness-to-pay for  $B$ , is more effective.
- ▶ When  $\delta > 1$ , consumers are more heterogeneous in  $Q$  than in  $B$ .
- ▶ Pricing by quantity becomes more effective.

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<sup>5</sup>In fact  $\eta$  and  $\gamma$  also have impacts on the heterogeneity. As the impacts are somewhat less apparent, we do not discuss them here.

## ADSL vs. 3G/4G

- ▶ Does our theory apply to the current practices?
- ▶ Currently, few data services are priced by minutes.
  - ▶ Supply side: Controlling the quantity is more direct than controlling time usage.
  - ▶ Consumer side: Pricing by minutes is not revenue-maximizing.
- ▶ ADSL is typically priced by bandwidth.
  - ▶ ADSL consumers are **more heterogeneous** in applications they prefer (and thus in **bandwidth**).
  - ▶ Therefore, pricing by bandwidth is more effective.
- ▶ 3G/4G is typically priced by quantity.
  - ▶ Few 3G/4G consumers use speed-demanding applications. Most of them spend most of the time on simple browsing/searching. They are **less heterogeneous in bandwidth**.
  - ▶ Pricing by quantity is thus more effective.

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## Extensions

- ▶ The model may be further extended in the following ways:
  - ▶ General utility functions:  $U(B, M, \theta) = U^Q(Q, \theta) + U^B(B, \theta)$ .
  - ▶ Bandwidth-insensitive utility functions:  $U(B, M, \theta) = U(Q, \theta)$ .
  - ▶ Aggregate bandwidth costs.
  - ▶ Disutility of waiting.
- ▶ In the presence of the last two supply-side issues:
  - ▶ Pricing by minutes is still suboptimal.
  - ▶ Pricing by bandwidth becomes relatively more attractive.

## Conclusions

- ▶ Three pricing metrics for data services are studied.
  - ▶ Pricing by minutes, bandwidth, or quantity.
- ▶ Either pricing by bandwidth or pricing by quantity can be optimal.
  - ▶ Pricing by minutes is the worst in mitigating **information asymmetry**. The remaining moral hazard problem is the most significant.
  - ▶ Whether the seller should price by bandwidth or quantity also depends on the effectiveness of mitigating information asymmetry.
- ▶ Why is information asymmetry critical?
  - ▶ We want to earn revenues at the consumer side.
  - ▶ We do not know how consumers **like** our product.
  - ▶ We do not know how consumers will **use** our product.
- ▶ After-sales selections are also important when we design returns, warranties, and many other consumer-related policies.