

Block Ciphers and DES

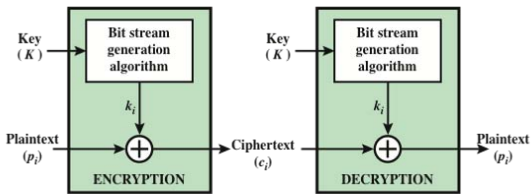
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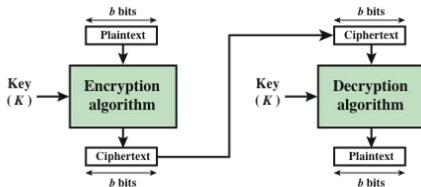
Block vs. Stream Ciphers

- 🌐 **Stream Cipher:** encrypt data one bit or one byte at a time.
 - ☀️ RC4 (classic examples: Vigenère cipher, Vernam cipher)
- 🌐 **Block Cipher:** encrypt data one block at a time; the size of a block ranges from several bytes to few tens of bytes.
 - ☀️ DES (Data Encryption Standard)
 - ☀️ AES (Advanced Encryption Standard)
- 🌐 Using certain modes of operation (to be discussed later), a block cipher can be used to achieve the effect of a stream cipher.
- 🌐 Block ciphers have a broader range of applications.

Block vs. Stream Ciphers (cont.)



(a) Stream Cipher Using Algorithmic Bit Stream Generator



(b) Block Cipher

Source: Figure 3.1, Stallings 2014

Block Ciphers

- 🌐 n-bit blocks of plaintext \longrightarrow n-bit blocks of ciphertext
- 🌐 $2^n!$ possible reversible mappings
- 🌐 Similar to classical substitution cipher when the block size is small
- 🌐 Arbitrary reversible mappings not practical for a large block size. The mapping itself is the key, requiring $n \times 2^n$ bits; e.g., $64 \times 2^{64} = 2^{70} \approx 10^{21}$.
- 🌐 **Approximation** (such as the Feistel Cipher) needed

Block Substitution

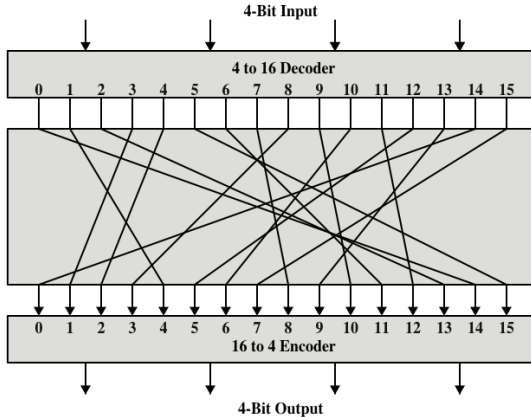


Figure 3.2 General n -bit- n -bit Block Substitution (shown with $n = 4$)

Source: Figure 3.2, Stallings 2014

Substitution Cipher

Plaintext	Ciphertext
0000	1110
0001	0100
0010	1101
0011	0001
0100	0010
0101	1111
0110	1011
0111	1000
1000	0011
1001	1010
1010	0110
1011	1100
1100	0101
1101	1001
1110	0000
1111	0111

Ciphertext	Plaintext
0000	1110
0001	0011
0010	0100
0011	1000
0100	0001
0101	1100
0110	1010
0111	1111
1000	0111
1001	1101
1010	1001
1011	0110
1100	1011
1101	0010
1110	0000
1111	0101

Source: Table 3.1, Stallings 2014

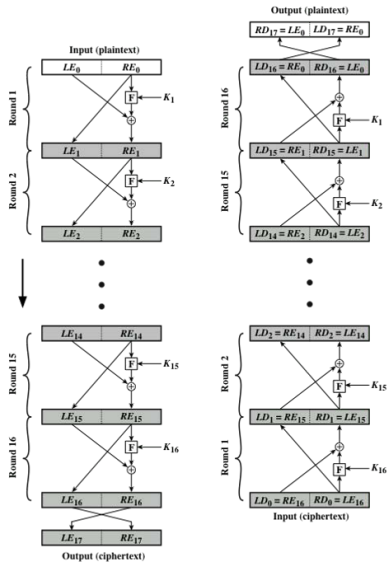
Diffusion and Confusion

- 🌐 Basic building blocks for any cipher
- 🌐 Methods for thwarting statistical cryptanalysis
- 🌐 Ideally, all statistics of the ciphertext should be independent of the particular key used.
- 🌐 **Diffusion**: make the statistical relationship between the plaintext and ciphertext as complex as possible, by having each plaintext digit (bit) affect the value of many ciphertext digits (bits).
Achieved by permutation plus some function.
- 🌐 **Confusion**: make the statistical relationship between the ciphertext and the encryption key as complex as possible.
Achieved by complex substitution (linear substitutions such as the Caesar cipher would add little confusion).

The Feistel Cipher

- 🌐 Approximate the ideal substitution cipher by a product cipher, i.e., two or more basic ciphers in sequence
- 🌐 Alternate **substitutions** and **permutations** (cf. confusion and diffusion)
- 🌐 Its structure used by many symmetric block ciphers currently in use

Feistel Encryption and Decryption



Source: Figure 3.3, Stallings 2014

Parameters of Feistel Networks

- 🌐 Block size: typically 64 bits
- 🌐 Key size: typically 128 bits
- 🌐 Number of rounds: typically 16 rounds
- 🌐 Subkey generation algorithm: the more complex the more secure
- 🌐 Round function: the more complex the more secure

Correctness of Feistel Decryption

Let's just check if the last round of decryption is correct.

Input: $LD_0 = RE_{16}$ and $RD_0 = LE_{16}$.

On the encryption side,

$$\begin{aligned}LE_{16} &= RE_{15} \\ RE_{16} &= LE_{15} \oplus F(RE_{15}, K_{16})\end{aligned}$$

On the decryption side,

$$\begin{aligned}LD_1 &= RD_0 = LE_{16} = RE_{15} \\ RD_1 &= LD_0 \oplus F(RD_0, K_{16}) \\ &= RE_{16} \oplus F(RE_{15}, K_{16}) \\ &= [LE_{15} \oplus F(RE_{15}, K_{16})] \oplus F(RE_{15}, K_{16}) \\ &= LE_{15}\end{aligned}$$

Feistel Example

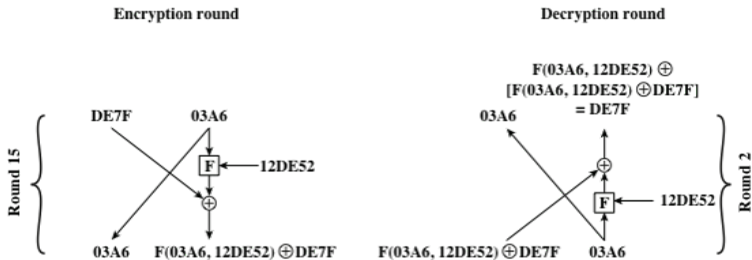


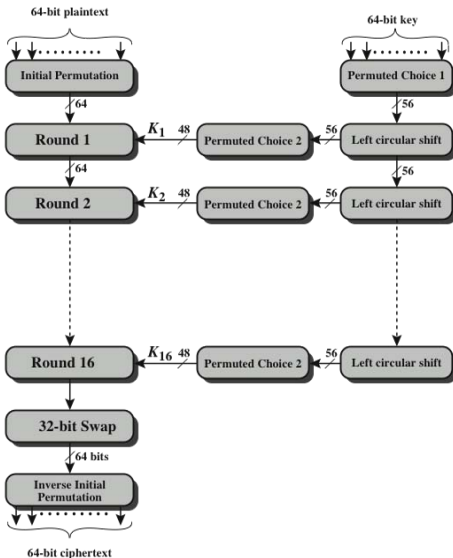
Figure 3.4 Feistel Example

Source: Figure 3.4, Stallings 2014

The Data Encryption Standard (DES)

- 🌐 Most widely used (and most important)
- 🌐 Data encrypted in **64-bit blocks** using a **56-bit key**
- 🌐 History:
 - ☀ LUCIFER (1971): developed by IBM; 128-bit keys.
 - ☀ Refined LUCIFER (1973): developed by IBM and NSA; 56-bit keys (**to fit on a chip**).
 - ☀ DES (1977): proposed by IBM and adopted by NBS, now NIST; a slightly modified version of the Refined LUCIFER; 56-bit keys.

DES Encryption



Source: Figure 3.5, Stallings 2014

Permutation Tables for DES

(a) Initial Permutation (IP)

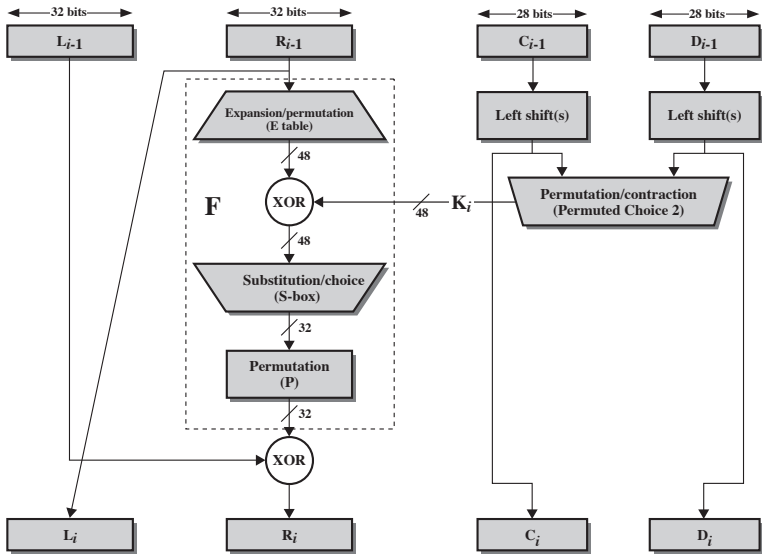
58	50	42	34	26	18	10	2
60	52	44	36	28	20	12	4
62	54	46	38	30	22	14	6
64	56	48	40	32	24	16	8
57	49	41	33	25	17	9	1
59	51	43	35	27	19	11	3
61	53	45	37	29	21	13	5
63	55	47	39	31	23	15	7

(b) Inverse Initial Permutation (IP^{-1})

40	8	48	16	56	24	64	32
39	7	47	15	55	23	63	31
38	6	46	14	54	22	62	30
37	5	45	13	53	21	61	29
36	4	44	12	52	20	60	28
35	3	43	11	51	19	59	27
34	2	42	10	50	18	58	26
33	1	41	9	49	17	57	25

Source: Table 3.2, Stallings 2010

Single Round of DES



Source: Figure 3.6, Stallings 2010

Permutation Tables for DES (cont.)

(c) Expansion Permutation (E)

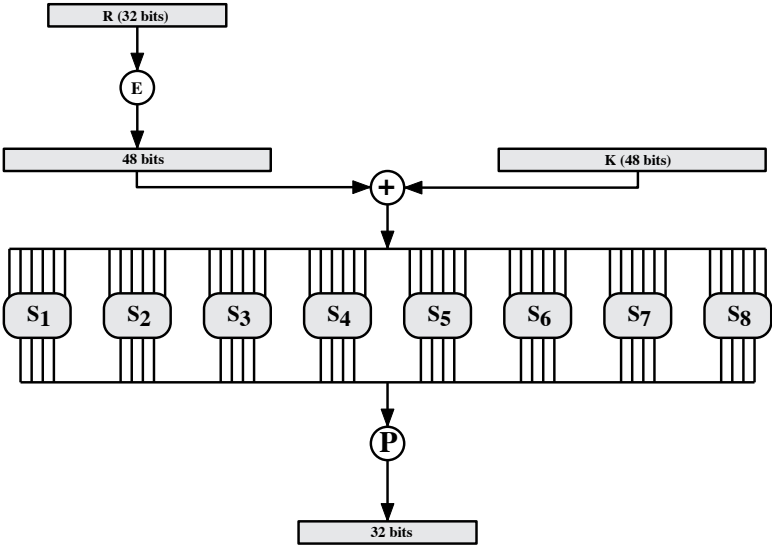
32	1	2	3	4	5
4	5	6	7	8	9
8	9	10	11	12	13
12	13	14	15	16	17
16	17	18	19	20	21
20	21	22	23	24	25
24	25	26	27	28	29
28	29	30	31	32	1

(d) Permutation Function (P)

16	7	20	21	29	12	28	17
1	15	23	26	5	18	31	10
2	8	24	14	32	27	3	9
19	13	30	6	22	11	4	25

Source: Table 3.2, Stallings 2010

Calculation of F(R,K)



Source: Figure 3.7, Stallings 2010

DES S-Boxes

S_1

14	4	13	1	2	15	11	8	3	10	6	12	5	9	0	7
0	15	7	4	14	2	13	1	10	6	12	11	9	5	3	8
4	1	14	8	13	6	2	11	15	12	9	7	3	10	5	0
15	12	8	2	4	9	1	7	5	11	3	14	10	0	6	13

S_2

15	1	8	14	6	11	3	4	9	7	2	13	12	0	5	10
3	13	4	7	15	2	8	14	12	0	1	10	6	9	11	5
0	14	7	11	10	4	13	1	5	8	12	6	9	3	2	15
13	8	10	1	3	15	4	2	11	6	7	12	0	5	14	9

S_3

10	0	9	14	6	3	15	5	1	13	12	7	11	4	2	8
13	7	0	9	3	4	6	10	2	8	5	14	12	11	15	1
13	6	4	9	8	15	3	0	11	1	2	12	5	10	14	7
1	10	13	0	6	9	8	7	4	15	14	3	11	5	2	12

S_4

7	13	14	3	0	6	9	10	1	2	8	5	11	12	4	15
13	8	11	5	6	15	0	3	4	7	2	12	1	10	14	9
10	6	9	0	12	11	7	13	15	1	3	14	5	2	8	4
3	15	0	6	10	1	13	8	9	4	5	11	12	7	2	14

S_5

2	12	4	1	7	10	11	6	8	5	3	15	13	0	14	9
14	11	2	12	4	7	13	1	5	0	15	10	3	9	8	6
4	2	1	11	10	13	7	8	15	9	12	5	6	3	0	14
11	8	12	7	1	14	2	13	6	15	0	9	10	4	5	3

S_6

12	1	10	15	9	2	6	8	0	13	3	4	14	7	5	11
10	15	4	2	7	12	9	5	6	1	13	14	0	11	3	8
9	14	15	5	2	8	12	3	7	0	4	10	1	13	11	6
4	3	2	12	9	5	15	10	11	14	1	7	6	0	8	13

S_7

4	11	2	14	15	0	8	13	3	12	9	7	5	10	6	1
13	0	11	7	4	9	1	10	14	3	5	12	2	15	8	6
1	4	11	13	12	3	7	14	10	15	6	8	0	5	9	2
6	11	13	8	1	4	10	7	9	5	0	15	14	2	3	12

S_8

13	2	8	4	6	15	11	1	10	9	3	14	5	0	12	7
1	15	13	8	10	3	7	4	12	5	6	11	0	14	9	2
7	11	4	1	9	12	14	2	0	6	10	13	15	3	5	8
2	1	14	7	4	10	8	13	15	12	9	0	3	5	6	11

Source: Table 3.3, Stallings 2010

(a) Input Key

1	2	3	4	5	6	7	8
9	10	11	12	13	14	15	16
17	18	19	20	21	22	23	24
25	26	27	28	29	30	31	32
33	34	35	36	37	38	39	40
41	42	43	44	45	46	47	48
49	50	51	52	53	54	55	56
57	58	59	60	61	62	63	64

Source: Table 3.4, Stallings 2010

DES Key Scheduling (cont.)

(b) Permuted Choice One (PC-1)

57	49	41	33	25	17	9
1	58	50	42	34	26	18
10	2	59	51	43	35	27
19	11	3	60	52	44	36
63	55	47	39	31	23	15
7	62	54	46	38	30	22
14	6	61	53	45	37	29
21	13	5	28	20	12	4

(c) Permuted Choice Two (PC-2)

14	17	11	24	1	5	3	28
15	6	21	10	23	19	12	4
26	8	16	7	27	20	13	2
41	52	31	37	47	55	30	40
51	45	33	48	44	49	39	56
34	53	46	42	50	36	29	32

(d) Schedule of Left Shifts

Round Number	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
Bits Rotated	1	1	2	2	2	2	2	2	1	2	2	2	2	2	2	1

Source: Table 3.4, Stallings 2010

The Avalanche Effect

- 🌐 The **avalanche effect** of an encryption algorithm refers to the property that **a small change in either the plaintext or the key should produce a significant change in the ciphertext.**
- 🌐 In particular, a change in one bit of the plaintext or the key should produce a change in many bits of the ciphertext.
- 🌐 DES exhibits a strong avalanche effect.

Avalanche Effect in DES

(a) Change in Plaintext		(b) Change in Key	
Round	Number of bits that differ	Round	Number of bits that differ
0	1	0	0
1	6	1	2
2	21	2	14
3	35	3	28
4	39	4	32
5	34	5	30
6	32	6	32
7	31	7	35
8	29	8	34
9	42	9	40
10	44	10	38
11	32	11	31
12	30	12	33
13	30	13	28
14	26	14	26
15	29	15	34
16	34	16	35

Source: Table 3.5, Stallings 2006

Breaking DES

- 🌐 Searching the 56-bit key space (about 7.2×10^{16} keys)
 - ☀️ The “DES Cracker” machine by the Electronic Frontier Foundation (1998)
 - 👷 It cost less than \$250,000 to build
 - 👷 The attack took less than three days
- 🌐 Exploiting the characteristics of DES
 - ☀️ Weakness in the S-boxes (none found so far)
 - ☀️ Timing attacks (unlikely to succeed)
(Timing attacks exploit the fact that an encryption or decryption implementation may take different amounts of time on different inputs.)






Exhaustive Key Search

Key size (bits)	Cipher	Number of Alternative Keys	Time Required at 10^9 decryptions/s	Time Required at 10^{13} decryptions/s
56	DES	$2^{56} \approx 7.2 \times 10^{16}$	2^{55} ns = 1.125 years	1 hour
128	AES	$2^{128} \approx 3.4 \times 10^{38}$	2^{127} ns = 5.3×10^{21} years	5.3×10^{17} years
168	Triple DES	$2^{168} \approx 3.7 \times 10^{50}$	2^{167} ns = 5.8×10^{33} years	5.8×10^{29} years
192	AES	$2^{192} \approx 6.3 \times 10^{57}$	2^{191} ns = 9.8×10^{40} years	9.8×10^{36} years
256	AES	$2^{256} \approx 1.2 \times 10^{77}$	2^{255} ns = 1.8×10^{60} years	1.8×10^{56} years
26 characters (permutation)	Monoalphabetic	$26! = 4 \times 10^{26}$	2×10^{26} ns = 6.3×10^9 years	6.3×10^6 years

(Note: the exponents are distorted, due to graphics incompatibility.)

Source: Table 3.5, Stallings 2014

Block Cipher Design Principles

-  **Number of rounds:** chosen so that known cryptanalytic efforts require greater effort than a brute-force key search attack.
-  **Design of function F**
 -  Design criteria for F: nonlinearity, strict avalanche, bit independence, etc.
 -  S-Box design: any change to the input vector to an S-box results in random-looking changes to the output.
In DES, S-Boxes are the main component of the function F and the only non-linear part of the entire DES.
-  **Key schedule algorithm:** maximize the difficulty of deducing individual subkeys and of working back to the main key.

Function F and S-Box Design Criteria

Nonlinearity

Strict avalanche:

Any output bit should change with probability $\frac{1}{2}$ when any single input bit is inverted.

Guaranteed avalanche:

Guaranteed avalanche of order γ ensures that, for a 1-bit input change, at least γ output bits change.

Bit independence:

Any two output bits should change independently when any single input bit is inverted.

S-Box Design Approaches

- 🌐 Random with or without testing
- 🌐 Man-made: manual with simple supporting math
- 🌐 Math-made: applying math principles to achieve proven security